

2017-2018 ACM-ICPC, NEERC, Moscow Subregional Contest

比赛情况

题号	A	B	C	D	E	F	G	H	I	J
状态	O	-	O	O	-	O	O	O	O	-

O 在比赛中通过 Ø 赛后通过! 尝试了但是失败了 - 没有尝试

比赛时间

2020-06-14 13:00-18:00

提交记录

A: Wrong Answer 2020-06-14 13:47 *3

A: Accepted 2020-06-14 13:49

D: Accepted 2020-06-14 14:00

H: Wrong Answer 2020-06-14 15:19 *3

H: Accepted 2020-06-14 15:28

C: Accepted 2020-06-14 15:42

G: Accepted 2020-06-14 16:16

F: Wrong Answer 2020-06-14 17:46 *7

F: Accepted 2020-06-14 17:58

题解

A.Advertising Strategy

目标订购者为 n 总花费 k a_i 为当前订购者数量, 每天可以花费 x_i 使得 $b_i = a_i + x_i$ 之后 $a_{i+1} = b_i + \min(b_i, \lfloor \frac{n-b_i}{2} \rfloor)$ 问达到目标订购的最少天数

题解: 显然至少要有花费要在最后一天给出, 其他钱应该是早花早受益, 于是在第一天和最后一天花钱。枚举在第一天花的钱, 模拟天数增加直到剩下的人可以一次性花钱补全。

code:

```
#include<bits/stdc++.h>
```

```
#define ll long long
#define pii_pair<int,int>
#define mp_make_pair
#define pb push_back
#define fi first
#define se second
#define rep(i,a,b) for(int i=(a);i<=(b);i++)
#define per(i,a,b) for(int i=(a);i>=(b);i--)
#define show1(a) cout<<#a<<" = "<<a<<endl
#define show2(a,b) cout<<#a<<" = "<<a<<" ; "<<#b<<" = "<<b<<endl
using namespace std;
const ll INF = 1LL<<60;
const int inf = 1<<30;
const int maxn = 2e5+5;
inline void fastio() {ios::sync_with_stdio(false);cin.tie(0);cout.tie(0);}

ll n,k;

int main()
{
    fastio();
    cin>>n>>k;
    ll res=INF;
    for(ll i=1;i<k;i++)
    {
        ll now = k-i,ans = 1;
        while(now<n-i){
            now = now + min(now,(n-now)/2);
            ans++;
        }
        res = min(res,ans);
    }
    cout<<res<<endl;
    return 0;
}
```

C.Carpet

要将 $n(1 \leq n \leq 100000)$ 个点的树铺在 1000000×20 的坐标平面里，使得边与边无端点之外的重合点。

题解：树剖，将轻儿子往上放，重儿子往右放。

code:

```
#include<bits/stdc++.h>
using namespace std;
```

```
const int N=1e5+5;
int n,head[N],cnt,sz[N];
int x[N],y[N],tot[25];
struct Node
{
    int nxt,to;
}Edges[N*2];
void addedge(int u,int v)
{
    Edges[++cnt].nxt=head[u];
    head[u]=cnt;
    Edges[cnt].to=v;
}
void dfs1(int u,int f)
{
    sz[u]=1;
    for(int i=head[u];~i;i=Edges[i].nxt)
    {
        int v=Edges[i].to;
        if(v==f)continue;
        dfs1(v,u),sz[u]+=sz[v];
    }
}
void dfs2(int u,int f,int d)
{
    y[u]=d,x[u]=++tot[d];
    int maxv=0,k=0;
    for(int i=head[u];~i;i=Edges[i].nxt)
    {
        int v=Edges[i].to;
        if(v==f)continue;
        if(sz[v]>maxv)maxv=sz[v],k=v;
    }
    for(int i=head[u];~i;i=Edges[i].nxt)
    {
        int v=Edges[i].to;
        if(v==f||v==k)continue;
        dfs2(v,u,d+1);
    }
    if(k)dfs2(k,u,d);
}
int main()
{
    memset(head,-1,sizeof(head));
    scanf("%d",&n);
    for(int i=1;i<n;i++)
    {
        int u,v;
        scanf("%d%d",&u,&v);
        addedge(u,v),addedge(v,u);
    }
}
```

```
dfs1(1,0),dfs2(1,0,1);  
for(int i=1;i<=n;i++)printf("%d %d\n",x[i],y[i]);  
return 0;  
}
```

D. Decoding of Varints

题意

语文题，就是给了一个类似128进制的定义，然后给一个未知的序列，对于这个序列中的每一个数字，如果这个数字 ≥ 0 的时候，这个数就变为这个数字的二倍，奇遇的时候是这个数字相反数的二倍再减一，把新数字搞成他之前定义的形式中的每一位，现在给你每一位，让你装换回去。

数据范围

$n \leq 10000$

题解

就 模拟啊

代码

```
#include <stdio.h>  
#include <string.h>  
#include <stdlib.h>  
#include <algorithm>  
using namespace std;  
const int N = 1e4+5;  
int a[N];  
typedef unsigned long long ll;  
int main()  
{  
    int n;  
    scanf("%d",&n);  
    for (int i = 1;i<= n;i++)scanf("%d",&a[i]);  
    for (int i = 1;i<= n;i++) {  
        int y = i;  
        ll ans = 0,x = 1;  
        while (y <= n) {  
            if (a[y]>=128) {  
                ans += (a[y]-128)*x;  
                x *= 128;  
                y++;  
            } else {  
                ans -= (a[y]+127)*x;  
                x *= 128;  
                y++;  
            }  
        }  
        printf("%d ",ans);  
    }  
}
```

```

        ans += a[y]*x;
        i = y;
        break;
    }
}
if (ans & 1)
    printf("%lld\n", (long long)0 - ((ans - 1) / 2 + 1));
else
    printf("%llu\n", ans / 2);
}
return 0;
}

```

F. Fake or Leak

反正就是大模拟。。码力和英语都好差，写了好多次才过

```

#include<bits/stdc++.h>
#define ll long long
#define pii_ pair<int,int>
#define mp_ make_pair
#define pb push_back
#define fi first
#define se second
#define rep(i,a,b) for(int i=(a);i<=(b);i++)
#define per(i,a,b) for(int i=(a);i>=(b);i--)
#define show1(a) cout<<#a<<" = "<<a<<endl
#define show2(a,b) cout<<#a<<" = "<<a<<"; "<<#b<<" = "<<b<<endl
using namespace std;
const ll INF = 1LL<<60;
const int inf = 1<<30;
const int maxn = 1005;
inline void fastio() {ios::sync_with_stdio(false);cin.tie(0);cout.tie(0);}

int n,m,k,a,t;
int
nowTime[maxn],nowSolve[maxn],has[maxn],bestTime[maxn],in[maxn],finalAc[maxn]
];
char op;
string name,lastName[maxn],rn[maxn];
map<string,int> mp;

void GG()
{
    cout<<"Fake"<<endl;
    exit(0);
}

```

```
struct node
{
    string name;
    int solve,time,fac;
    bool operator <(node e){
        if(solve!=e.solve) return solve>e.solve;
        if(time!=e.time) return time<e.time;
        if(fac!=e.fac) return fac<e.fac;
        return name<e.name;
    }
};
vector<node> team;
vector<string> res;
int main()
{
    fastio();
    cin>>n>>m>>k;
    rep(j,1,m){
        cin>>name;
        mp[name] = j;
        rn[j] = name;
        rep(i,1,n){
            cin>>op>>a>>t;
            if(op=='+'){
                nowTime[j] += 20*(a-1)+t;
                nowSolve[j]++;
                bestTime[j] += 20*(a-1)+t;
                finalAc[j] = max(finalAc[j],t);
            }
            else bestTime[j] += 20*a+240;
        }
    }
    rep(j,1,k){
        cin>>laterName[j];
        name = laterName[j];
        int ti = 0,so = 0,fac=0;
        has[mp[name]] = 1;
        rep(i,1,n){
            cin>>op>>a>>t;
            if(op=='+'){
                ti += 20*(a-1)+t;
                so++;
                fac = max(fac,t);
            }
        }
        team.pb((node){name,so,ti,fac});
    }
    rep(i,1,m){
        if(!has[i]){
            team.pb((node){rn[i],nowSolve[i],nowTime[i],finalAc[i]});
        }
    }
}
```

```

        if(nowSolve[i]<n) team.pb((node){rn[i],n,bestTime[i],240});
        else team.pb((node){rn[i],n,bestTime[i],finalAc[i]});
    }
}
sort(team.begin(),team.end());
int cnt = 0,flag = 0;
for(node x:team){
    if(has[mp[x.name]]) flag = 1;
    if(flag && cnt<k){
        if(has[mp[x.name]]) {res.pb(x.name);cnt++;}
        else {
            in[mp[x.name]]++;
            if(in[mp[x.name]]==2) GG();
        }
    }
}
rep(i,1,k){
    if(res[i-1]!=laterName[i]) GG();
}
cout<<"Leaked"<<endl;
return 0;
}

```

G. God of Winds

一个网格图，每个小正方形的边有个权值。可以任选格子，加入一个顺时针或者逆时针的风，用来改变格子周围四条边的权值。右和下是正方形，题目给定权值，问这样的图是否存在。注意到每条边只有两个格子能影响它，假设每个 (i,j) 处的格子有 $x_{i,j}$ 个顺时针方向的风，那么整个图就是个等式的差分约束，对每个变量连边然后 dfs 一遍，判断会不会有矛盾的点。

```

#include<bits/stdc++.h>
#define ll long long
#define pii_ pair<int,int>
#define mp_ make_pair
#define pb push_back
#define fi first
#define se second
#define rep(i,a,b) for(int i=(a);i<=(b);i++)
#define per(i,a,b) for(int i=(a);i>=(b);i--)
#define show1(a) cout<<#a<<" = "<<a<<endl
#define show2(a,b) cout<<#a<<" = "<<a<<" "; cout<<#b<<" = "<<b<<endl
using namespace std;
const ll INF = 1LL<<60;
const int inf = 1<<30;
inline void fastio() {ios::sync_with_stdio(false);cin.tie(0);cout.tie(0);}
const int maxn = 250005;
int n,m;

```

```
int r[505][505],c[505][505];
int vis[maxn],head[maxn],tot;
ll dis[maxn];
struct edge
{
    int u,v,w,nxt;
}es[maxn<<2];
void add(int u,int v,int w)
{
    es[++tot] = (edge){u,v,w,head[u]};
    head[u] = tot;
}
void dfs(int u)
{
    vis[u] = 1;
    for(int i=head[u];~i;i=es[i].nxt){
        int v = es[i].v,w = es[i].w;
        if(!vis[v]){
            dis[v] = dis[u] + w;
            dfs(v);
        }else{
            if(dis[v]!=dis[u]+w){
                cout<<"No"<<endl;
                exit(0);
            }
        }
    }
}
int main()
{
    fastio();
    memset(head,-1,sizeof(head));
    cin>>n>>m;
    rep(i,0,n-1)rep(j,0,m-1){
        cin>>r[i][j]>>c[i][j];
    }
    rep(i,0,n-1){
        rep(j,0,m-1){
            int u = i*m+j,v = (i-1+n)%n*m + j;
            add(u,v,r[i][j]);add(v,u,-r[i][j]);
            int x = i*m+j,y = i*m + (j-1+m)%m;
            add(x,y,-c[i][j]);add(y,x,c[i][j]);
        }
    }
    dfs(0);
    cout<<"Yes"<<endl;
    return 0;
}
```


H. Hilarious Cooking

题意

希望你构造一个序列，序列相邻两个数差值不超过1，和为 T 然后其中有一些给定位置的给定数字，问能否构造成功。

数据范围

序列的长度 $n \leq 2 \times 10^9$ $1 \leq T \leq 10^{18}$

给定的数字个数 $1 \leq m \leq 100000$

题解

可以考虑，构造这样一个序列可以保证在一个范围内连续，因为除非在最大值和最小值的位置，其他情况下都可以找到一个位置+1或者找到一个位置-1并符合序列规定，那么我们要做的就是确定这个序列可能的最大值和可能的最小值。注意到其限制其实就是给定的数字，最大值在两个给定的数字构造一个像山顶或山坡的形状，最小值就构造一个山谷或山坡一样的形状即可，要注意一些细节，比如最开始和最后的。

代码

```
#include <stdio.h>
#include <string.h>
#include <stdlib.h>
#include <algorithm>
using namespace std;
typedef long long ll;
const int N = 1e5+5;
ll x[N],y[N];
ll T,n;
ll calc(ll x,ll y) {
    if (x > y || y < 0) return 0;
    ll nx = x < 0 ? -x : 0;
    return (x+y)*(y-x+1)/2 + (nx+1)*nx/2;
}
int main()
{
    scanf("%lld",&T);
    ll minn = 0;
    ll maxx = 0;
    int m;
    scanf("%lld%d",&n,&m);
    for (int i = 1;i <= m;i++) {
        scanf("%lld%lld",&x[i],&y[i]);
        minn += y[i];
        maxx += y[i];
    }
```

```
}
for (int i = 1; i < m; i++) {
    ll dis = x[i+1]-x[i];
    ll val = abs(y[i+1]-y[i]);
    if (val > dis) {
        printf("No\n");
        return 0;
    }
    if (dis == 1) { continue; }
    ll maxx = max(y[i+1],y[i]) + (dis-val)/2;
    ll minx = min(y[i+1],y[i]) - (dis-val)/2;
    if (val % 2 == dis % 2) {
        maxn += calc(max(y[i+1],y[i]),maxx) +
        calc(min(y[i+1],y[i]),maxx) - max(maxx,0ll);
        minn += calc(minx,max(y[i+1],y[i])) +
        calc(minx,min(y[i+1],y[i])) - max(minx,0ll);
    } else {
        maxn += calc(max(y[i+1],y[i]),maxx) +
        calc(min(y[i+1],y[i]),maxx);
        minn += calc(minx,max(y[i+1],y[i])) +
        calc(minx,min(y[i+1],y[i]));
    }
    maxn -= y[i+1]+y[i];
    minn -= y[i+1]+y[i];
}
maxn += calc(y[1]+1,y[1]+x[1]-1);
minn += calc(y[1]-x[1]+1,y[1]-1);
maxn += calc(y[m]+1,y[m]+n-x[m]);
minn += calc(y[m]-n+x[m],y[m]-1);
if (T<= maxn && T >= minn) { printf("Yes\n"); }
else { printf("No\n"); }
return 0;
}
```

replay

比赛总结

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